

# A Padé-based factorization-free algorithm for identifying the eigenvalues missed by a generalized symmetric eigensolver

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## SUMMARY

When computing the solution of a generalized symmetric eigenvalue problem of the form  $\mathbf{K}\mathbf{u} = \lambda\mathbf{M}\mathbf{u}$ , the Sturm sequence check, also known as the inertia check, is the most popular method for reporting the number of missed eigenvalues within a range  $[\sigma_L, \sigma_R]$ . This method requires the factorization of the matrices  $\mathbf{K} - \sigma_L\mathbf{M}$  and  $\mathbf{K} - \sigma_R\mathbf{M}$ . When the size of the problem is reasonable and the matrices  $\mathbf{K}$  and  $\mathbf{M}$  are assembled, these factorizations are possible. When the eigensolver is equipped with an iterative solver, which is nowadays the preferred choice for large-scale problems, the factorization of  $\mathbf{K} - \sigma\mathbf{M}$  is not desired or feasible and therefore the inertia check cannot be performed. To this effect, the purpose of this paper is to present a factorization-free algorithm for detecting and identifying the eigenvalues that were missed by an eigensolver equipped with an iterative linear equation solver within an interval of interest  $[\sigma_L, \sigma_R]$ . This algorithm constructs a scalar, rational, transfer function whose poles are exactly the eigenvalues of the symmetric pencil  $(\mathbf{K}, \mathbf{M})$ , approximates it by a Padé expansion, and computes the poles of this approximation to detect and identify the missed eigenvalues. The proposed algorithm is illustrated with an academic numerical example. Its potential for real engineering applications is also demonstrated with a large-scale structural vibrations problem. Copyright © 2009 John Wiley & Sons, Ltd.

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## 1. INTRODUCTION

The numerical solution of a generalized large-scale eigenvalue problem of the form

$$\mathbf{K}\mathbf{u} = \lambda\mathbf{M}\mathbf{u} \quad (1)$$

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where  $\mathbf{K}$  and  $\mathbf{M}$  are real symmetric positive definite matrices of size  $n$ , which arises in many engineering problems. For example, in structural dynamics, the circular frequencies  $\lambda_k = \omega_k^2$  and corresponding eigenmodes  $\mathbf{u}_k$  are typically determined to prevent resonant excitations, design a controller, or simply build a reduced-order computational model to predict the transient linear response of the structure. Confidence in any of these applications requires, among others, verifying that all relevant eigenpairs have been computed by the eigensolver.

A corollary of Sylvester's inertia theorem [1] implies that the number of eigenvalues lying in a real interval  $[\sigma_L, \sigma_R]$  can be determined by factoring the matrices  $\mathbf{K} - \sigma_R \mathbf{M}$  and  $\mathbf{K} - \sigma_L \mathbf{M}$ , counting their number of negative pivots, and computing the difference between these two numbers. Frequently, the number of negative eigenvalues of  $\mathbf{K} - \sigma \mathbf{M}$  is called the Sturm sequence number in engineering references [2, p. 239]. For large-scale three-dimensional problems with ten million or more degrees of freedom, factoring the aforementioned matrices can be computationally prohibitive. In this case, experts often re-apply their eigensolvers in the space that is  $\mathbf{M}$ -orthogonal to the previously computed eigenmodes in order to identify any missed eigenpair. This idea was recently exploited in [3] to implement an iterative scheme for computing *all*  $p$  lowest eigenpairs of a symmetric positive definite pencil  $(\mathbf{K}, \mathbf{M})$ , for a specified  $p$ . This scheme relies on the monotonic convergence of the eigenvalues sequenced in the increasing order. As such, it can be applied to identify the eigenvalues missed in the range  $[\sigma_L, \sigma_R]$  only if  $\sigma_L < \lambda_1$ , where  $\lambda_1$  denotes the smallest eigenvalue of the symmetric positive definite pencil  $(\mathbf{K}, \mathbf{M})$ . Unfortunately, this can be restrictive.

To the best of the authors' knowledge, no post-processing technique that does not require the factorization of a matrix related to  $\mathbf{K}$  and/or  $\mathbf{M}$  is currently available for checking whether an eigensolver applied to the solution of problem (1) has missed some eigenvalues in an arbitrary range of interest  $[\sigma_L, \sigma_R]$ . Hence, the main objective of this paper is to fill this gap in the computational literature. We emphasize that our intent is to provide a technique to identify missed modes (if any) when factoring the aforementioned matrices which is computationally prohibitive. This constraint eliminates the possibility of using the Sturm sequence because no factorization is available.

As  $\mathbf{K} - \lambda \mathbf{M}$  is singular only when  $\lambda$  is an eigenvalue of  $(\mathbf{K}, \mathbf{M})$ , it follows that the poles of a scalar rational function such as  $\mathbf{l}^T (\mathbf{K} - \lambda \mathbf{M})^{-1} \mathbf{r}$ , where  $\mathbf{l}$  and  $\mathbf{r}$  are two real vectors and the superscript T designates the transpose, are the eigenvalues of  $(\mathbf{K}, \mathbf{M})$ . It has been shown that such a rational function can be accurately and efficiently approximated by an appropriate Padé expansion [4–9]. The numerical algorithm proposed in this paper for identifying the eigenvalues missed in a given range of interest builds on the aforementioned observations. It is described and illustrated in the remainder of this paper which is organized as follows.

In Section 2, the overall approach proposed in this paper for identifying the missed eigenvalues is presented and justified. A key component of this approach is a scalar rational function that is approximated by a Padé expansion whose computation is detailed in Section 3. The proposed method is laid out in Section 4. It is illustrated in Section 5 with several examples, including a large-scale problem from structural dynamics. Its performance is also assessed and its potential for real engineering applications is highlighted. Finally, conclusions are offered in Section 6.

## 2. OVERALL APPROACH AND NUMERICAL ALGORITHM

Consider the following scalar transfer function:

$$H(\sigma) = \mathbf{b}^T (\mathbf{K} - \sigma \mathbf{M})^{-1} \mathbf{b} \quad (2)$$

where  $\mathbf{K}$  and  $\mathbf{M}$  are two symmetric positive definite matrices, and  $\mathbf{b}$  is an arbitrary vector. The poles of  $H$  are exactly the eigenvalues of the generalized symmetric eigenvalue problem

$$(\mathbf{K} - \lambda\mathbf{M})\mathbf{u} = \mathbf{0}$$

Hence, checking the eigenvalues of  $(\mathbf{K}, \mathbf{M})$  is equivalent to checking the poles of  $H$ . Therefore, the approach proposed in this paper for identifying the eigenvalues missed by an eigensolver in an interval of interest  $[\sigma_L, \sigma_R]$  consists in first computing a good approximation of  $H$ , then computing the poles of this approximation, and finally analyzing these poles to identify the sought-after missed eigenvalues. To this effect, it is first shown that  $H$  is a rational function.

### 2.1. Rational expression of $H$

Since  $\mathbf{K}$  and  $\mathbf{M}$  are symmetric, the pencil  $(\mathbf{K}, \mathbf{M})$  is diagonalizable. The diagonalization of this pencil can be written as

$$\mathbf{K}\mathbf{U} = \mathbf{M}\mathbf{U}\mathbf{\Lambda}$$

where  $\mathbf{U}^T\mathbf{M}\mathbf{U} = \mathbf{I}$  and  $\mathbf{\Lambda}$  is a diagonal matrix whose entries are the eigenvalues of  $(\mathbf{K}, \mathbf{M})$ . The columns of the matrix  $\mathbf{U}$  are the corresponding eigenvectors.

Observing that  $\mathbf{U}^{-1} = \mathbf{U}^T\mathbf{M}$ , one obtains

$$\mathbf{K} = \mathbf{M}\mathbf{U}\mathbf{\Lambda}\mathbf{U}^T\mathbf{M}$$

and

$$\mathbf{K} - \sigma\mathbf{M} = \mathbf{M}\mathbf{U}(\mathbf{\Lambda} - \sigma\mathbf{I})\mathbf{U}^T\mathbf{M}$$

It follows that

$$(\mathbf{K} - \sigma\mathbf{M})^{-1} = \mathbf{U}(\mathbf{\Lambda} - \sigma\mathbf{I})^{-1}\mathbf{U}^T$$

and therefore

$$H(\sigma) = \mathbf{b}^T\mathbf{U}(\mathbf{\Lambda} - \sigma\mathbf{I})^{-1}\mathbf{U}^T\mathbf{b} \quad (3)$$

The above result can be rewritten as

$$H(\sigma) = \sum_{k=1}^n \frac{(\mathbf{U}^T\mathbf{b})_k^2}{\lambda_k - \sigma} \quad (4)$$

where  $(\cdot)_k$  denotes the  $k$ th component of a vector and  $\lambda_k$  is the  $k$ th diagonal entry in  $\mathbf{\Lambda}$ . This shows that  $H$  is a rational function of  $\sigma$  with a numerator that is a polynomial of degree at most equal to  $n - 1$  and a denominator of degree at most equal to  $n$ , where  $n$  denotes the size of the matrix  $\mathbf{K}$ . Equation (4) also emphasizes that the poles of  $H$  are exactly the eigenvalues of the pencil  $(\mathbf{K}, \mathbf{M})$ .

### 2.2. One-point Padé approximation for local accuracy

A rational function such as  $H$  is typically well represented by a Padé approximant. The poles of this approximant can be expected to converge toward the poles of  $H$ , which are exactly the eigenvalues of the pencil  $(\mathbf{K}, \mathbf{M})$ . A Padé approximation of type  $(l/m)$  of the function  $H$  around

$\sigma_0$ , an arbitrary real number for which however the matrix  $\mathbf{K} - \sigma_0\mathbf{M}$  is nonsingular, is defined as a rational function of the form

$$H_{l,m}(\sigma) = \frac{p_0 + p_1(\sigma - \sigma_0) + \dots + p_l(\sigma - \sigma_0)^l}{1 + q_1(\sigma - \sigma_0) + \dots + q_m(\sigma - \sigma_0)^m} \quad (5)$$

whose Taylor expansion around  $\sigma_0$  matches the first  $l + m + 1$  terms of the Taylor expansion of  $H$  around the same point—that is,

$$H(\sigma) = H_{l,m}(\sigma) + \mathcal{O}((\sigma - \sigma_0)^{l+m+1})$$

As the numerator and denominator of  $H$  are polynomials of degrees equal at most to  $n - 1$  and  $n$ , respectively, a Padé approximation of type  $((m - 1)/m)$  is considered here and simply denoted by  $H_m$ . The coefficients of the polynomials are computed so that the approximant matches the first  $2m$  terms of the following Taylor expansion around  $\sigma_0$ ,

$$H_m(\sigma_0) = H(\sigma_0) \quad \text{and} \quad \frac{d^j H_m}{d\sigma^j}(\sigma_0) = \frac{d^j H}{d\sigma^j}(\sigma_0) \quad \forall j < 2m$$

For the sake of completeness, the derivatives of  $H$  at  $\sigma_0$  are computed. The transfer function  $H$  introduced in Equation (2) can be rewritten as follows:

$$\begin{aligned} H(\sigma) &= \mathbf{b}^T [\mathbf{K} - \sigma_0\mathbf{M} - (\sigma - \sigma_0)\mathbf{M}]^{-1} \mathbf{b} \\ &= \mathbf{b}^T \{ (\mathbf{K} - \sigma_0\mathbf{M}) [\mathbf{I} - (\sigma - \sigma_0)(\mathbf{K} - \sigma_0\mathbf{M})^{-1}\mathbf{M}] \}^{-1} \mathbf{b} \\ &= \mathbf{b}^T [\mathbf{I} - (\sigma - \sigma_0)(\mathbf{K} - \sigma_0\mathbf{M})^{-1}\mathbf{M}]^{-1} (\mathbf{K} - \sigma_0\mathbf{M})^{-1} \mathbf{b} \end{aligned} \quad (6)$$

The Taylor expansion around  $\sigma_0$  is

$$H(\sigma) = \sum_{q=0}^{+\infty} (\sigma - \sigma_0)^q \mathbf{b}^T [(\mathbf{K} - \sigma_0\mathbf{M})^{-1}\mathbf{M}]^q (\mathbf{K} - \sigma_0\mathbf{M})^{-1} \mathbf{b} \quad (7)$$

Since

$$H(\sigma_0) = \mathbf{b}^T (\mathbf{K} - \sigma_0\mathbf{M})^{-1} \mathbf{b}$$

and

$$\frac{d^j H}{d\sigma^j}(\sigma_0) = j! \mathbf{b}^T [(\mathbf{K} - \sigma_0\mathbf{M})^{-1}\mathbf{M}]^j (\mathbf{K} - \sigma_0\mathbf{M})^{-1} \mathbf{b} \quad \text{for } j < 2m$$

it follows that

$$H_m(\sigma_0) = \mathbf{b}^T (\mathbf{K} - \sigma_0\mathbf{M})^{-1} \mathbf{b} \quad (8)$$

and

$$\frac{d^j H_m}{d\sigma^j}(\sigma_0) = j! \mathbf{b}^T [(\mathbf{K} - \sigma_0\mathbf{M})^{-1}\mathbf{M}]^j (\mathbf{K} - \sigma_0\mathbf{M})^{-1} \mathbf{b} \quad \text{for } j < 2m \quad (9)$$

2.3. Multi-point Padé approximation for global accuracy

The Padé approximant  $H_m$  can be expected to converge first locally around the expansion point  $\sigma_0$ . However, the polynomial degree  $m$  can be expected to grow large before  $H_m$  becomes a globally accurate approximation of  $H$ . As the main focus here is on checking for missing eigenvalues of  $(\mathbf{K}, \mathbf{M})$  in a range of interest  $[\sigma_L, \sigma_R]$ , a global approximation of  $H$  over the entire interval  $[\sigma_L, \sigma_R]$  is preferable. Such a global approximation can be provided by the multi-point Padé method which matches the moments of  $H$  at  $I$  discrete points  $(\sigma_i)_{1 \leq i \leq I} \in [\sigma_L, \sigma_R]$ . If one chooses to fit the value of  $H$  and its  $2J - 1$  derivatives at each point  $\sigma_i$ , the resulting multi-point Padé approximation of  $H$  satisfies

$$\forall i \in \{1, \dots, I\}, \quad H_m(\sigma_i) = H(\sigma_i) \quad \text{and} \quad \frac{d^j H_m}{d\sigma^j}(\sigma_i) = \frac{d^j H}{d\sigma^j}(\sigma_i) \quad \forall j < 2J$$

In this case, one can write

$$H(\sigma) = H_m(\sigma) + \mathcal{O}\left(\prod_{i=1}^I (\sigma - \sigma_i)^{2J}\right)$$

where  $m = IJ$ . The distribution of the  $I$  points  $\sigma_i$  within the interval  $[\sigma_L, \sigma_R]$  is arbitrary; it can be a uniform or nonuniform one. The multi-point Padé approximation can also consider a variable number of derivatives at the points  $\sigma_i$ . However, for the sake of simplicity, it is chosen here to consider the same number of derivatives at each point  $\sigma_i$ .

2.4. General form of the proposed algorithm

Let  $\mathbf{U}_c$  denote the matrix storing the  $n_c < n$  computed eigenvectors whose associated computed eigenvalues lie within the interval of interest  $[\sigma_L, \sigma_R]$ . If  $\mathbf{b}$  is a random vector, it can be assumed that

$$(\mathbf{U}_c^T \mathbf{b})_k \neq 0 \quad \forall k \leq n_c$$

On the other hand if  $\mathbf{b}$  is chosen in the null space of  $\mathbf{U}_c^T$ —that is, if

$$\mathbf{U}_c^T \mathbf{b} = \mathbf{0} \tag{10}$$

it follows from Equation (4) that

$$H(\sigma) = \sum_{k | \lambda_k \in \Lambda_{nc}} \frac{(\mathbf{U}^T \mathbf{b})_k^2}{\lambda_k - \sigma} \tag{11}$$

where  $\Lambda_{nc}$  is the union of the set of missed eigenvalues within  $[\sigma_L, \sigma_R]$  and the set of eigenvalues that are outside this interval. Hence, if  $\mathbf{b}$  is chosen to satisfy Equation (10),  $H$  identifies with the expression given in (11) and its poles become exactly the union of the eigenvalues that were missed in  $[\sigma_L, \sigma_R]$  and those that lie outside this interval. This result leads to the following algorithm for identifying the eigenvalues of the pencil  $(\mathbf{K}, \mathbf{M})$  that were missed by an eigensolver in the interval  $[\sigma_L, \sigma_R]$ .

1. Gather in  $\mathbf{U}_c$  the computed eigenvectors associated with the computed eigenvalues that lie in the interval  $[\sigma_L, \sigma_R]$ .
2. Generate a random vector  $\tilde{\mathbf{b}}$ .

3. Compute  $\mathbf{b}$ , the projection of  $\tilde{\mathbf{b}}$  onto the space orthogonal to  $\mathbf{U}_c$ ;  $\mathbf{b}$  satisfies  $\mathbf{U}_c^T \mathbf{b} = \mathbf{0}$ .
4. Compute  $H_m$ , a multi-point Padé approximation of  $H$  in  $[\sigma_L, \sigma_R]$ .
5. Identify the poles of  $H_m$  in the interval  $[\sigma_L, \sigma_R]$ .

Next, the computation of the multi-point Padé approximation  $H_m$  is discussed.

### 3. COMPUTATION OF THE MULTI-POINT PADÉ APPROXIMATION

#### 3.1. Relationship with Krylov subspaces

In [4], a nonsymmetric Lanczos process was developed for computing a single-point Padé approximant of a transfer function such as  $H$  and was shown to be computationally efficient. In [5], an approach based on the rational nonsymmetric Lanczos algorithm was proposed for computing a multi-point approximation of a rational function. Unfortunately, the nonsymmetric Lanczos algorithm can suffer a loss of orthogonality and may even break down. For this reason, a robust alternative based on a two-sided rational Arnoldi algorithm was proposed in [6, 10]. In the specific case of the rational function  $H$  given in (2) where the matrices are symmetric and the vectors are the transpose of each other, the two-sided Arnoldi algorithm reduces to a symmetric Lanczos algorithm.

All numerical methods outlined above for computing a Padé approximant of a rational function such as  $H$  project the matrices  $\mathbf{K}$  and  $\mathbf{M}$ , and the vector  $\mathbf{b}$ , onto a carefully constructed matrix  $\mathbf{V}$ . As such, they are often referred to as methods of ‘reduction by projection’. Furthermore, all of these methods rely on Krylov subspaces for constructing the appropriate matrix  $\mathbf{V}$ . Here, a Krylov subspace is denoted by

$$\mathcal{K}_q(\check{\mathbf{K}}, \check{\mathbf{b}}) = \text{span}\{\check{\mathbf{b}}, \check{\mathbf{K}}\check{\mathbf{b}}, \dots, \check{\mathbf{K}}^{q-1}\check{\mathbf{b}}\} \quad (12)$$

where  $\check{\mathbf{K}}$  is a matrix of interest and  $\check{\mathbf{b}}$  is a starting vector, and the projections of  $\mathbf{K}$ ,  $\mathbf{M}$ , and  $\mathbf{b}$  onto  $\mathbf{V}$  are denoted by

$$\hat{\mathbf{K}} = \mathbf{V}^T \mathbf{K} \mathbf{V}, \quad \hat{\mathbf{M}} = \mathbf{V}^T \mathbf{M} \mathbf{V} \quad \text{and} \quad \hat{\mathbf{b}} = \mathbf{V}^T \mathbf{b} \quad (13)$$

In [6, Theorem 3.1], it is proved that if the columns of the constructed  $\mathbf{V}$  are such that their span verifies the following property:

$$\text{span}\{\mathbf{V}\} \supset \bigoplus_{i=1}^J \mathcal{K}_J((\mathbf{K} - \sigma_i \mathbf{M})^{-1} \mathbf{M}, (\mathbf{K} - \sigma_i \mathbf{M})^{-1} \mathbf{b}) \quad (14)$$

the reduced matrices  $\hat{\mathbf{K}}$  and  $\hat{\mathbf{M}}$ , and the reduced vector  $\hat{\mathbf{b}}$  satisfy

$$\hat{\mathbf{b}}^T [(\hat{\mathbf{K}} - \sigma_i \hat{\mathbf{M}})^{-1} \hat{\mathbf{M}}]^j (\hat{\mathbf{K}} - \sigma_i \hat{\mathbf{M}})^{-1} \hat{\mathbf{b}} = \mathbf{b}^T [(\mathbf{K} - \sigma_i \mathbf{M})^{-1} \mathbf{M}]^j (\mathbf{K} - \sigma_i \mathbf{M})^{-1} \mathbf{b} \quad (15)$$

where  $j = 0, 1, \dots, 2J - 1$ . From Equations (8), (9), and (15) above, it follows that when the property (14) is satisfied, the function

$$H_m(\sigma) = \hat{\mathbf{b}}^T (\hat{\mathbf{K}} - \sigma \hat{\mathbf{M}})^{-1} \hat{\mathbf{b}} = (\mathbf{V}^T \mathbf{b})^T (\mathbf{V}^T \mathbf{K} \mathbf{V} - \sigma \mathbf{V}^T \mathbf{M} \mathbf{V})^{-1} \mathbf{V}^T \mathbf{b} \quad (16)$$

matches the values of  $H$  and its  $2J - 1$  derivatives at the points  $(\sigma_i)_{1 \leq i \leq J}$ , and therefore is a multi-point Padé approximant of  $H$ .

### 3.2. Lanczos algorithm for computing the subspace of projection

Algorithm 1 describes a symmetric Lanczos algorithm with full reorthogonalization for constructing the columns of a matrix  $\mathbf{V}$ , of dimension  $n \times IJ$ , that satisfies  $\mathbf{V}^T \mathbf{M} \mathbf{V} = \mathbf{I}$ ,  $\mathbf{U}_c^T \mathbf{M} \mathbf{V} = \mathbf{0}$ , and the property (14). After  $\mathbf{V}$  is built, the reduced-order quantities (13) can be constructed and the sought-after multi-point Padé approximant (16) of  $H$  can be computed.

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#### Algorithm 1 Symmetric Lanczos process

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(Initialize)  $p = 1$ 
for  $i = 1, 2, \dots, I$  do
  for  $j = 1, 2, \dots, J$  do
    if  $j == 1$  then
       $\mathbf{v} = (\mathbf{K} - \sigma_i \mathbf{M})^{-1} \mathbf{b}$ 
    else
       $\mathbf{v} = (\mathbf{K} - \sigma_i \mathbf{M})^{-1} (\mathbf{I} - \mathbf{U}_c (\mathbf{U}_c^T \mathbf{U}_c)^{-1} \mathbf{U}_c^T) \mathbf{M} \mathbf{v}_{p-1}$ 
    end if
    if  $p > 1$  then
      (Orthogonalize)  $\mathbf{v} = \mathbf{v} - \mathbf{V}_{p-1} \mathbf{V}_{p-1}^T \mathbf{M} \mathbf{v}$ 
    end if
     $\gamma = \sqrt{\mathbf{v}^T \mathbf{M} \mathbf{v}}$ 
    if  $\gamma == 0$  then
      Reset  $\mathbf{v}$  to a random vector
      (Orthogonalize)  $\mathbf{v} = \mathbf{v} - \mathbf{V}_{p-1} \mathbf{V}_{p-1}^T \mathbf{M} \mathbf{v}$ .
       $\gamma = \sqrt{\mathbf{v}^T \mathbf{M} \mathbf{v}}$ .
    end if
    (Update)  $\mathbf{v}_p = \mathbf{v} / \gamma$ 
     $p = p + 1$ 
  end for
end for

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Note that the construction of  $\mathbf{V}$  requires the solution of  $m = IJ$  linear systems ( $J$  linear systems with  $(\mathbf{K} - \sigma_i \mathbf{M})$  for  $I$  sample points  $\sigma_i$ ). When the norm  $\sqrt{\mathbf{v}^T \mathbf{M} \mathbf{v}}$  of a vector  $\mathbf{v}$  is zero, Algorithm 1 updates the columns of  $\mathbf{V}$  by a new random vector while maintaining the property (14). The reduced matrix  $\hat{\mathbf{M}}$  is replaced by the identity matrix. The reduced matrix  $\hat{\mathbf{K}}$  and the reduced vector  $\hat{\mathbf{b}}$  are built by direct projection.

## 4. MISSING EIGENVALUE FINDER

The discussion presented in the previous sections of this paper leads to proposing Algorithm 2 described below for the factorization-free identification of the eigenvalues missed in an arbitrary range of interest  $[\sigma_L, \sigma_R]$  by an eigensolver applied to the solution of the generalized symmetric eigenvalue problem (1).

Algorithm 2 is essentially a post-processing procedure. It does not depend on the eigensolver used to compute the eigenvectors  $\mathbf{U}_c$  and the corresponding eigenvalues. The inputs to this proposed algorithm are the computed eigenvectors  $\mathbf{U}_c$ , the range of interest  $[\sigma_L, \sigma_R]$ , the desired number of

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**Algorithm 2** Lanczos-based eigencheck

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**Require:** Computed  $\mathbf{M}$ -orthonormal eigenvectors  $\mathbf{U}_c$

**Require:** Range of interest  $[\sigma_L, \sigma_R]$

**Require:** Value for  $I$  (for sampling the above range of interest in  $I$  points  $(\sigma_i)_{1 \leq i \leq I}$ )

**Require:** Value for  $J$  (for matching at each point  $\sigma_i$  the first  $2J - 1$  derivatives of the scalar rational function  $H$ )

1. Sample  $[\sigma_L, \sigma_R]$  in  $I$  points  $(\sigma_i)_{1 \leq i \leq I}$
2. Generate a random vector  $\tilde{\mathbf{b}}$
3. Compute  $\mathbf{b} = (\mathbf{I} - \mathbf{U}_c(\mathbf{U}_c^T \mathbf{U}_c)^{-1} \mathbf{U}_c^T) \tilde{\mathbf{b}}$
4. Apply Algorithm 1 for constructing the matrix  $\mathbf{V}$  of dimension  $n \times m$  (where  $m = IJ$ ).
5. Compute the reduced matrices  $\hat{\mathbf{K}}$  and  $\hat{\mathbf{M}}$  (we recall  $\hat{\mathbf{M}} = \mathbf{I}$ )
6. Compute the *full* solution of the reduced-order  $m \times m$  generalized symmetric eigenvalue problem

$$\hat{\mathbf{K}}\mathbf{w} = \lambda \hat{\mathbf{M}}\mathbf{w} \tag{17}$$

7. Identify those computed eigenvalues of problem (17) which lie in the interval of interest  $[\sigma_L, \sigma_R]$
- 

sampling points  $I$ , and the integer  $J$  for which the Padé approximant of  $H$  matches its first  $2J - 1$  derivatives at the sampled points  $(\sigma_i)_{1 \leq i \leq I}$ . The projection of the randomly generated vector  $\tilde{\mathbf{b}}$  (Step 3) ensures that the vector  $\mathbf{b}$  satisfies the desired property

$$\mathbf{U}_c^T \mathbf{b} = \mathbf{0}$$

so that the poles of the computed Padé approximant correspond to the missed eigenvalues in the range of interest  $[\sigma_L, \sigma_R]$ .

Note that Algorithm 2 cannot guarantee that all of the missing modes will be identified by using any set of parameters  $(I, J)$ . But this trait is typical of many numerical algorithms. Even the numerical performance of the Sturm sequence check depends on a tolerance  $\varepsilon$ . When a pivot belongs to  $[-\varepsilon, \varepsilon]$ , the inertia check cannot determine whether the pivot is associated with a zero, positive, or negative eigenvalue. However, a *refinement* process can easily accompany Algorithm 2. Given a range  $[\sigma_L, \sigma_R]$  and  $I$  sampling points  $(\sigma_i)_{1 \leq i \leq I}$ , increasing  $J$  will improve the multi-point Padé approximation of  $H$  and the number of eigenvalues of problem (17) in the interval of interest  $[\sigma_L, \sigma_R]$  will converge to a fixed value. Given a range  $[\sigma_L, \sigma_R]$  and a value for  $J$ , increasing  $I$  will also improve the multi-point Padé approximation of  $H$ .

In Section 1, it was noted that experts often try to find the missed eigenpairs by re-applying their eigensolvers in the space that is  $\mathbf{M}$ -orthogonal to the previously computed eigenvectors. Here, it is concluded that this process is equivalent to Algorithm 2 with  $I = 1$  (a single sampling point) and therefore is unreliable for identifying all eigenvalues that have been missed in a specific range of interest  $[\sigma_L, \sigma_R]$ .

## 5. EXAMPLES

In this section, the proposed factorization-free Algorithm 2 for finding the eigenvalues missed in an interval of interest by an eigensolver is illustrated with two numerical examples. The convergence

for the number of eigenvalues of problem (17) in the interval of interest  $[\sigma_L, \sigma_R]$  is highlighted. The first example is of the academic type. It has the merit of being easily reproducible by the reader. It also demonstrates the main behavior and tendencies of the proposed Algorithm 2. The second example is associated with a large-scale structural vibrations problem. It highlights the potential of the proposed algorithm for real-engineering applications. In both cases, ARPACK [11] is chosen as the eigensolver.

5.1. Matrix BCSSTK16

Here, the generalized symmetric eigenvalue problem (1) is considered with  $\mathbf{K}$  set to the matrix BCSSTK16 from the Harwell–Boeing sparse matrix collection [12] whose size is  $n=4884$ , and  $\mathbf{M}$  set to the identity matrix. For reference, a complete solution of this problem is obtained using the MATLAB EIG routine. It is found that all 73 first eigenvalues of the pencil  $(\mathbf{K}, \mathbf{M})$  are equal to 1.0, the 74th one is equal to  $1.5 \times 10^5$ , and the largest eigenvalue is equal to  $4.9 \times 10^9$ .

To compute an approximation of the first 100 eigenpairs  $(\mathbf{U}_c, \Lambda_c)$  of the problem specified above, ARPACK is applied within MATLAB 7 in the shift–invert mode with a shift  $\sigma=0$ , a Krylov subspace of size equal to 200, and the convergence tolerance  $\tau=10^{-8}$ . For this purpose, the random starting vector is generated with the MATLAB commands

```
rand('state', 1);
resid = rand(size(K,1), 1);
```

Using these parameters, 67 of the 100 computed eigenpairs are found to have their eigenvalues in the range  $[0, 10]$ . As the problem considered here has exactly 73 eigenvalues in  $[0, 10]$ , this means that in this case, ARPACK has missed six eigenpairs in the range  $[0, 10]$ . The 67 computed eigenvalues in  $[0, 10]$  are equal to 1.0 up to machine precision.

Then, Algorithm 2 is applied to find the eigenvalues missed in the range  $[0, 10]$ . The random vector  $\mathbf{b}$  is generated with the MATLAB commands:

```
rand('state', 2);
btilde = rand(size(K, 1), 1);
```

The obtained results with  $I=1$  (one-point Padé approximation) are reported in Table I. More specifically, Table I reports for different matching points  $\sigma_0$ , the number of eigenvalues identified by Algorithm 2 as missed eigenvalues in  $[0, 10]$ , as a function of the polynomial degree  $m=J$

Table I. Problem BCSSTK16: number of identified missed eigenvalues in  $[0, 10]$  (specified after the + symbol).

	$\sigma_0=0$	$\sigma_0=5$	$\sigma_0=10$
ARPACK	67	67	67
ARPACK+poles of $H_2$	67+1	67+1	67+1
ARPACK+poles of $H_4$	67+2	67+2	67+1
ARPACK+poles of $H_8$	67+3	67+3	67+3
ARPACK+poles of $H_{16}$	67+6	67+5	67+5
ARPACK+poles of $H_{17}$	67+6	67+6	67+5
ARPACK+poles of $H_{18}$	67+6	67+6	67+6
ARPACK+poles of $H_{19}$	67+6	67+6	67+6
ARPACK+poles of $H_{20}$	67+6	67+6	67+6

Table II. Problem BCSSTK16: poles recovered by Algorithm 2 when  $\sigma_0=5$ .

Recovered values by Algorithm 2	
Poles of $H_2$	(1.0000, $5.0 \times 10^7$ )
Poles of $H_4$	(1.0000, 1.0917, $3.4 \times 10^7$ , $1.2 \times 10^8$ )
Poles of $H_6$	(1.0000, 1.0000, $3.0 \times 10^7$ , $4.5 \times 10^7$ , $8.7 \times 10^7$ , $3.7 \times 10^8$ )
Poles of $H_8$	(1.0000, 1.0000, 1.0000, $3.0 \times 10^7$ , $3.9 \times 10^7$ , $5.5 \times 10^7$ , ...)
Poles of $H_{12}$	(1.0000, 1.0000, 1.0000, 1.0000, $3.0 \times 10^7$ , $3.2 \times 10^7$ , ...)
Poles of $H_{16}$	(1.0000, 1.0000, 1.0000, 1.0000, 1.0000, $3.0 \times 10^7$ , ...)
Poles of $H_{17}$	(1.0000, 1.0000, 1.0000, 1.0000, 1.0000, 1.0716, ...)
Poles of $H_{18}$	(1.0000, 1.0000, 1.0000, 1.0000, 1.0000, 1.0000, ...)
Poles of $H_{19}$	(1.0000, 1.0000, 1.0000, 1.0000, 1.0000, 1.0000, ...)
Poles of $H_{20}$	(1.0000, 1.0000, 1.0000, 1.0000, 1.0000, 1.0000, ...)

characterizing the Padé approximant  $H_J$ . The reader can observe that once it has converged, Algorithm 2 retrieves all six eigenpairs that have been missed by ARPACK in the range  $[0, 10]$ . The convergence of this algorithm is shown however to depend on the matching point  $\sigma_0$ . Indeed, since all 73 eigenvalues lying in  $[0, 10]$  are equal to 1, the closer  $\sigma_0$  is to 1, the smaller can the degree  $m$  be expected.

To illustrate that Algorithm 2 recovers the correct eigenvalues in  $[0, 10]$ , Table II reports the smallest eigenvalues of the reduced pencil  $(\hat{\mathbf{K}}, \hat{\mathbf{M}})$  (up to the first six) when  $\sigma_0=5$ . The reader can observe that once it has identified six missed eigenvalues in the range  $[0,10]$  (with  $H_{17}$ ), Algorithm 2 refines the values for the eigenpairs that have been missed by ARPACK. However, Algorithm 2 can only guarantee that the recovered eigenvalues will be accurate up to  $\tau=10^{-8}$  as  $(\mathbf{U}_c, \mathbf{\Lambda}_c)$  were computed with that tolerance.

Next, a smaller tolerance  $\tau=10^{-16}$  is used and all other ARPACK parameters are kept the same. In this case, ARPACK finds all 73 eigenvalues that belong to the range  $[0, 10]$ , and the post-processing Algorithm 2 correctly finds that no eigenvalue has been missed in this range.

Finally, the post-processing Algorithm 2 is reapplied for finding the eigenvalues missed in the range  $[0, 10]$  using this time  $I>1$  matching points (multi-point Padé approximation). Table III reports for this case the minimal polynomial degree  $m^*=IJ^*$  (or for a given  $I$ , the minimum value of  $J, J^*$ ) that enables Algorithm 2 to identify six missed eigenvalues in the range  $[0, 10]$ , as a function of  $I$  and the choice of the sampling points  $(\sigma_i)_{1 \leq i \leq I}$ . Note that  $m^*$  is also the number of columns in  $\mathbf{V}$ . As expected from the global approximation property of the multi-point Padé algorithm,  $m^*$  is found to decrease with  $I$ .

### 5.2. Large-scale structural vibrations problem

Here, the generalized symmetric eigenvalue problem (1) is considered with  $\mathbf{K}$  and  $\mathbf{M}$  set to the stiffness and mass matrices, respectively, arising from the finite element discretization of the bolted joint model of Figure 1 whose size is  $n=322710$ . To construct an approximation of the first 100 eigenpairs  $(\mathbf{U}_c, \mathbf{\Lambda}_c)$  of this model, the P(arallel)ARPACK FORTRAN library functions PDSAUPD and PDSEUPD are applied within an in-house developed finite element structural analysis code in the shift–invert mode. The shift parameter is set to  $\sigma=0$ , the size of the Krylov subspace is set to 200, the convergence tolerance to  $\tau=10^{-16}$ , and the FETI-DP solver [13, 14] is chosen to solve all the resulting linear systems of equations. Furthermore, every 10th converged eigenpair

FINDING MISSING EIGENVALUES USING A PADÉ-BASED ALGORITHM

Table III. Problem BCSSTK16: smallest polynomial degree  $m = IJ$  for which Algorithm 2 identifies correctly all eigenvalues missed in  $[0, 10]$ .

I	$(\sigma_i)_{1 \leq i \leq I}$	$J^*$	$m^* = IJ^*$
1	{5}	17	17
2	{0, 10}	9	18
3	{0, 5, 10}	6	18
4	$\{0, \frac{10}{3}, \frac{20}{3}, 10\}$	4	16
5	{0, 2.5, 5, 7.5, 10}	3	15
6	$\{0, \frac{10}{6}, \frac{20}{6}, \frac{30}{6}, \frac{40}{6}, \frac{50}{6}, 10\}$	2	12

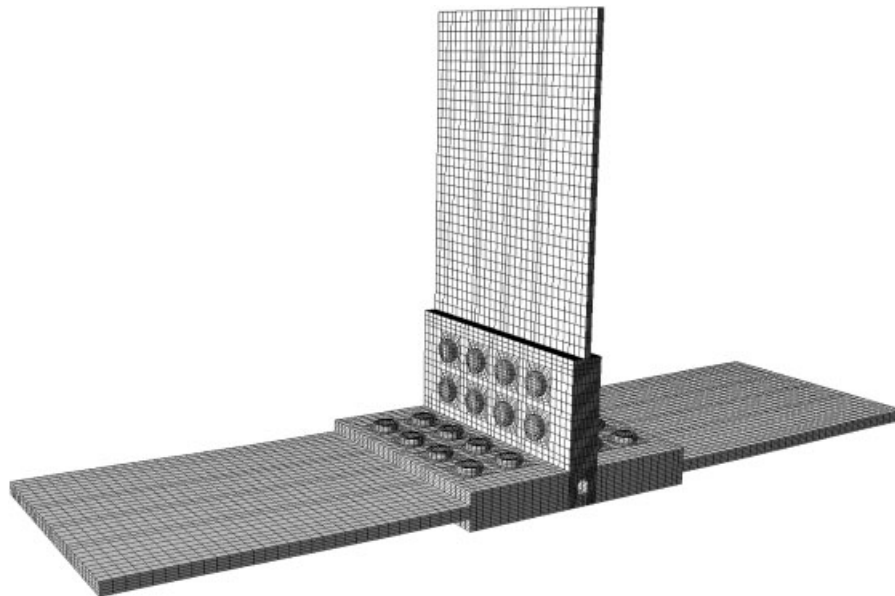


Figure 1. Large-scale finite element model of a bolted joint for modal analysis.

is deliberately excluded from the approximation  $(\mathbf{U}_c, \Lambda_c)$  so that some (or all) of the missed eigenvalues can be known *a priori* for the purpose of verification.

Using the above parameters, 92 of the 100 computed eigenpairs are found to have their eigenvalues in the range  $[0, 3.9478 \times 10^9]$  (corresponding to the frequency range  $[0, 1 \times 10^4]$  Hz) of which nine are deliberately excluded from the approximation. The eigenvalues of the excluded eigenpairs are  $3.4939 \times 10^7$ ,  $1.7928 \times 10^8$ ,  $3.3891 \times 10^8$ ,  $8.0245 \times 10^8$ ,  $1.0672 \times 10^9$ ,  $1.4863 \times 10^9$ ,  $2.4172 \times 10^9$ ,  $2.8610 \times 10^9$ , and  $3.7337 \times 10^9$ . The number of genuinely missed eigenvalues is unknown, but is most likely zero since the problem did not exhibit any clustering phenomenon and a tight convergence tolerance was chosen.

Next, Algorithm 2 is applied to find the excluded eigenvalues as well as those potentially missed in the range  $[0, 3.9478418 \times 10^9]$ . In Algorithm 1, the linear systems are solved iteratively by

Table IV. Large-scale structural vibrations problem: number of identified missed eigenvalues in  $[0, 3.9478 \times 10^9]$  (specified after the + symbol).

	$\sigma_0=0$	$\sigma_0=1.9739 \times 10^9$	$\sigma_0=3.9478 \times 10^9$
PARPACK <sup>†</sup>	83	83	83
PARPACK <sup>†</sup> + poles of $H_4$	83+2	83+2	83+1
PARPACK <sup>†</sup> + poles of $H_8$	83+5	83+4	83+2
PARPACK <sup>†</sup> + poles of $H_{12}$	83+7	83+6	83+3
PARPACK <sup>†</sup> + poles of $H_{16}$	83+7	83+8	83+4
PARPACK <sup>†</sup> + poles of $H_{18}$	83+8	83+9	83+5
PARPACK <sup>†</sup> + poles of $H_{20}$	83+8	83+9	83+5
PARPACK <sup>†</sup> + poles of $H_{22}$	83+9	83+9	83+6
PARPACK <sup>†</sup> + poles of $H_{24}$	83+9	83+9	83+6
PARPACK <sup>†</sup> + poles of $H_{26}$	83+9	83+9	83+7
PARPACK <sup>†</sup> + poles of $H_{28}$	83+9	83+9	83+7
PARPACK <sup>†</sup> + poles of $H_{30}$	83+9	83+9	83+8
PARPACK <sup>†</sup> + poles of $H_{32}$	83+9	83+9	83+8
PARPACK <sup>†</sup> + poles of $H_{34}$	83+9	83+9	83+8
PARPACK <sup>†</sup> + poles of $H_{36}$	83+9	83+9	83+9
PARPACK <sup>†</sup> + poles of $H_{38}$	83+9	83+9	83+9
PARPACK <sup>†</sup> + poles of $H_{40}$	83+9	83+9	83+9

PARPACK<sup>†</sup> denotes PARPACK results with selected eigenpairs excluded.

the FETI-DPH solver [15], a domain-decomposition-based iterative solver designed for indefinite problems of the form  $(\mathbf{K} - \sigma\mathbf{M})\mathbf{x} = \mathbf{f}$ . The results obtained with  $I = 1$  (one-point Padé approximation) are summarized in Table IV. More specifically, Table IV reports for different matching points  $\sigma_0$ , the number of eigenvalues identified by Algorithm 2 as missed eigenvalues in  $[0, 3.9478 \times 10^9]$ , as a function of the polynomial degree  $m = J$  characterizing the Padé approximant  $H_J$ . The reader can observe that once it has converged, Algorithm 2 retrieves all nine eigenpairs excluded from the approximation in the range  $[0, 3.9478 \times 10^9]$  and does not identify any other missed eigenvalue. Since all 92 eigenvalues lying in  $[0, 3.9478 \times 10^9]$  are evenly spaced, a value of  $\sigma_0$  in the middle of the range can be expected, and indeed is found, to be the best one.

To illustrate that Algorithm 2 recovers the correct eigenvalues in  $[0, 3.9478 \times 10^9]$ , Table V reports the smallest eigenvalues of the reduced pencil  $(\hat{\mathbf{K}}, \hat{\mathbf{M}})$  (up to the first nine ones) when  $\sigma_0 = 1.9739 \times 10^9$ . The reader can observe that once it has identified nine missed eigenvalues in the range  $[0, 3.9478 \times 10^9]$ , Algorithm 2 refines the actual values for the identified missed eigenpairs. Note that nine poles for  $H_{20}$  and  $H_{22}$  are in the range  $[0, 3.9478 \times 10^9]$  and that the smallest pole is not exact for  $H_{20}$  but more accurate for  $H_{22}$ .

Finally, the post-processing Algorithm 2 is reapplied for finding the eigenvalues missed in the range  $[0, 3.9478 \times 10^9]$  using this time  $I > 1$  matching points (multi-point Padé approximation). Table VI reports for this case the minimal polynomial degree  $m^* = IJ^*$  (or for a given  $I$ , the minimum value of  $J, J^*$ ) that enables Algorithm 2 to identify nine missed eigenvalues in the range  $[0, 3.9478 \times 10^9]$ , as a function of  $I$  and the choice of the sampling points  $(\sigma_i)_{1 \leq i \leq I}$ . Note that  $m^*$  is also the dimension of the reduced pencil  $(\hat{\mathbf{K}}, \hat{\mathbf{M}})$  and the total number of linear solves

FINDING MISSING EIGENVALUES USING A PADÉ-BASED ALGORITHM

Table V. Large-scale structural vibrations problem: poles recovered by Algorithm 2 when  $\sigma_0 = 1.9739 \times 10^9$ .

Recovered values by Algorithm 2	
Poles of $H_4$	$(2.2807 \times 10^9, 2.4173 \times 10^9)$
Poles of $H_8$	$(9.6607 \times 10^8, 1.4863 \times 10^9, 2.4172 \times 10^9, 2.8452 \times 10^9)$
Poles of $H_{12}$	$(3.1697 \times 10^8, 8.0244 \times 10^8, 1.0672 \times 10^9, 1.4863 \times 10^9, 2.4172 \times 10^9, 2.8610 \times 10^9)$
Poles of $H_{16}$	$(1.7875 \times 10^8, 3.3889 \times 10^9, 8.0245 \times 10^9, 1.0672 \times 10^9, 1.4863 \times 10^9, 2.4172 \times 10^9, 2.8610 \times 10^9, 3.7349 \times 10^9)$
Poles of $H_{20}$	$(3.4946 \times 10^7, 1.7928 \times 10^8, 3.3891 \times 10^8, 8.0245 \times 10^8, 1.0672 \times 10^9, 1.4863 \times 10^9, 2.4172 \times 10^9, 2.8610 \times 10^9, 3.7337 \times 10^9, \dots)$
Poles of $H_{22}$	$(3.4939 \times 10^7, 1.7928 \times 10^8, 3.3891 \times 10^8, 8.0245 \times 10^8, 1.0672 \times 10^9, 1.4863 \times 10^9, 2.4172 \times 10^9, 2.8610 \times 10^9, 3.7337 \times 10^9, \dots)$
Poles of $H_{24}$	$(3.4939 \times 10^7, 1.7928 \times 10^8, 3.3891 \times 10^8, 8.0245 \times 10^8, 1.0672 \times 10^9, 1.4863 \times 10^9, 2.4172 \times 10^9, 2.8610 \times 10^9, 3.7337 \times 10^9, \dots)$
Poles of $H_{26}$	$(3.4939 \times 10^7, 1.7928 \times 10^8, 3.3891 \times 10^8, 8.0245 \times 10^8, 1.0672 \times 10^9, 1.4863 \times 10^9, 2.4172 \times 10^9, 2.8610 \times 10^9, 3.7337 \times 10^9, \dots)$
Poles of $H_{28}$	$(3.4939 \times 10^7, 1.7928 \times 10^8, 3.3891 \times 10^8, 8.0245 \times 10^8, 1.0672 \times 10^9, 1.4863 \times 10^9, 2.4172 \times 10^9, 2.8610 \times 10^9, 3.7337 \times 10^9, \dots)$

Table VI. Large-scale structural vibrations problem: smallest polynomial degree  $m = IJ$  for which Algorithm 2 identifies correctly all eigenvalues missed in  $[0, 3.9478 \times 10^9]$ .

I	$(\sigma_i)_{1 \leq i \leq I}$	$J^*$	$m^* = IJ^*$
1	$\{1.9739 \times 10^9\}$	18	18
2	$\{0, 3.9478 \times 10^9\}$	9	18
3	$\{0, 1.9739 \times 10^9, 3.9478 \times 10^9\}$	6	18
4	$\{0, 1.3159 \times 10^9, 2.6319 \times 10^9, 3.9478 \times 10^9\}$	4	16
5	$\{0, 9.8696 \times 10^8, 1.9739 \times 10^9, 2.9609 \times 10^9, 3.9478 \times 10^9\}$	3	15
6	$\{0, 7.8956 \times 10^8, 1.5791 \times 10^9, 2.3687 \times 10^9, 3.1582 \times 10^9, 3.9478 \times 10^9\}$	3	18

performed in Algorithm 1. As expected from the global approximation property of the multi-point Padé algorithm,  $J^*$  and  $m^*$  are found to decrease with  $I$ .

6. CONCLUSIONS

To the best of the authors' knowledge, a factorization-free algorithm for detecting and identifying the eigenvalues missed by an eigensolver equipped with an iterative linear equation solver within an interval of interest  $[\sigma_L, \sigma_R]$  has been presented for the first time. This algorithm constructs a scalar, rational, transfer function whose poles are exactly the eigenvalues of the symmetric pencil  $(\mathbf{K}, \mathbf{M})$ , approximates it by a Padé expansion, and computes the poles of this approximation to detect and identify the missed eigenvalues. The proposed algorithm was illustrated with an academic numerical example. Its potential for real engineering applications was also demonstrated with a large-scale structural vibrations problem.

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